Document Numbers: J16/97-0042
WG21/N1080
Date: June 5, 1997
Reply To: Bill Gibbons
bill@gibbons.org

# More Proposed Resolutions for Core-III Issues

## Introduction

Between the March 1997 and July 1997 meetings, several members of Core-III met informally to discuss open template issues as well as a few other open core-III issues. The following people participated in the discussions:

Mike Ball
Bill Gibbons
Josee Lajoie
John Spicer
Daveed Vandevoorde
Tom Wilcox
John Wilkinson

The group came to agreement on almost all of the open Core-III issues. In this paper we propose changes (or lack of change) to the working paper to resolve the issues as described. The detailed working paper changes have not been done, but should be available at the July 1997 meeting.

Most of the issues are from N1065 and/or N1053; the descriptions of those issues are repeated here for clarity. The group also agreed with the "semi-editorial" issue resolutions in N1065 section 3, as repeated below, with an additional point of clarification on 3.16

Thanks to Sun for hosting the informal meeting, and to the participants - especially Josée Lajoie and John Spicer, who had to travel from out of the area to attend.

## Proposed Resolutions

From N1065, the Core-III list of issues compiled at Nashua

(1.2) [Core issue 765] 14.2 [temp.names] ¶4
"The syntax does not allow the keyword 'template' in '`typename expr.template C<parm>`'"
It is not clear whether the template keyword is needed in the above example. If so, the grammar must be changed to allow it. This issue should be investigated.

Change the grammar to allow the forms:

```
A::B::template C<5>::x
typename A::B::template C<5>::Y
```

The first would otherwise be impossible to parse within a template when nontype template arguments are used. The second could be parsed because C cannot be a nontype (because of the typename keyword), but the "template" keyword should be required anyway.

(1.3) [Core issue 736, parts 2 & 3] 14.6 [temp.res]
"How can/must typename be used?"

(2) The working paper does not require that typename be applied only to qualified names. Should there be such a restriction?
(3) The working paper does not require that typename be applied only to dependent names. Should there be such a restriction?

> Keep status quo: the typename keyword may only be applied to qualified names, but those names need not be dependent. The typename keyword may only be used in templates (including the return type of a function template).

(1.4) [N1053 issue 6.49 and issue 6.53 item 3]
"Where are partial specializations allowed?"
It is not clear whether a partial specialization must be declared in the class or namespace of which it is a member. There are cases for member templates where such a rule would prevent specialization entirely. The restrictions, if any, should be explicitly stated in the working paper.

> Allow partial specializations of class templates (but not function templates) to be initially declared outside the class or namespace in which the master class template is declared.

(1.5) [N1053 issue 6.51]
"Clarification of nontype dependency rules in partial specializations"
The restrictions in 14.5.4 [temp.class.spec] ¶5 item 2 make a large class of partial specializations ill-formed for no apparent reason. The restriction should probably be relaxed, possibly by restoring the less restrictive wording from a previous version of the working paper.

> Restore the original, less severe restrictions on nontype template arguments as described in N1053 issues 6.51

(1.6) [N1053 issue 6.52]
"Clarification of ordering rules for nontype arguments in partial specializations."
The partial ordering rules for class template partial specializations are too restrictive with respect to nontype template parameters. The rules should be reformulated to allow additional obviously correct orderings.

> Apply partial ordering rules to nontype template arguments as described in N1053 issue 6.52

(1.7) [N1053 issue 6.55]
"Interaction of partial ordering with default arguments and ellipsis parameters"
The working paper does not give clear rules for the handling of default arguments and ellipsis parameters when determining the partial ordering of function templates.

> The presence of unused ellipsis parameters and default arguments has no effect on the partial ordering of function templates, as described in N1053 issue 6.55

(1.8) [N1053 issue 6.56]
"In which contexts should partial ordering of function templates be performed?"
In addition to overload resolution, there are additional contexts in which partial ordering of function templates could be used to resolve ambiguities between function template instances with identical function parameters (and possibly identical template arguments) but generated from different partial specializations:

- Taking the address of a template function instance
- Matching a declaration of an instance with a particular partial specialization (for friend declarations, explicit specialization and explicit instantiation)

- Selecting a placement delete function that matches a placement new operation.
  It might be useful to apply the partial ordering rules in these contexts.

  Apply partial ordering of function templates when they are used certain contexts other than function calls, e.g. taking the address of a function. Also apply partial ordering to find the match for an explicit (complete) specialization declaration, even though this context could be considered a "declaration" instead of "use". (On this last part, Bill Gibbons opposed the proposed resolution.)

(1.9) [Editorial box #6] 14.5.4 [temp.class.spec] ¶5
The restrictions on partial specialization based on the dependency of arguments on other arguments are too severe. The restrictions should be relaxed where possible.

  See 1.5

(1.10) [Editorial box #11] 14.6.4.1 [temp.point]
The rules describing the point of instantiation for function templates may be overly complex. Consideration should be given to simplifying them.

  When an external declaration uses template in a way that causes a point of instantiation, the point of instantiation is immediately after the external declaration. The change is to remove 14.6.4.1 paragraph 2. (There is no problem with default arguments; the point of instantiation must not be before the external declaration because of indirect recursive calls.)

(1.11) [Editorial box #12] 14.6.4.2 [temp.dep.candidate]
This section says that if the visibility of candidate functions with external linkage in additional translations units affects the meaning of the program, the behavior is undefined. The possibility of extending the rule to include candidate functions without external linkage should be considered.

  No change. The presence of functions with internal linkage in other parts a program has no effect on the correctness of a template instantiation, even if making such functions visible would have changed the semantics of the program.

(1.12) [Public comment #23] 14.6.5 ¶2 and 3.4.2
The example does not match the argument-dependent name lookup rules for friends stated in 3.4.2. The rules in 3.4.2 do not match those presented to the committee when the extended argument-dependent name lookup rules were added.

  The rules for type-dependent name lookup should be fixed to handle friend functions in the manner in which was originally agreed. This requires defining "associated classes" as well as "associated namespaces", and saying that friend functions of associated classes are visible in their respective namespaces during type-dependent lookup. Although this was agreed to and written down at one time, it was either not completely edited into the working paper or parts of it were lost.

  I recently posted this explanation to the core reflector:

  Friend function declarations may introduce a new declaration into an enclosing scope, but the *name* itself is not introduced into that scope (much like "local externs"). Such declarations are checked against any other declarations in that scope for consistency, but they are not found during an ordinary name lookup. If a subsequent declaration in that scope matches (i.e. it is a redeclaration), the name becomes visible in the normal manner.

During type-dependent lookup, any class which is considered when determining the set of "associated namespaces" becomes an "associated class". For the purposes of the current lookup only, all friends of associated classes become visible in their namespaces as if they were directly declared there.

(1.13) [Public comment #26 ¶9]
It is not clear whether a rethrow creates a new exception which shares the exception object with the old exception, or whether the result of the rethrow is the old exception itself. If it is the latter, then the state of the exception should probably change from “caught” to “uncaught” as a result of the rethrow. This issue is not discussed in the working paper.

A rethrow reactivates the caught exception; it does not create a new one. The state of "uncaught_exception" is updated to indicate that the exception is once again uncaught.

(1.14) [new issue in March]
[Core issue 737]
In the following example:

```
template<class T> struct A {
   typedef int B;
   A<T>::B b;
};
```

is the lookup of B considered dependent? If so, is the example ill-formed? In what contexts is the use of a qualifier to look in the current template a special case not subject to the usual dependent type restrictions? Under what circumstances is a base class member found using a derived class qualifier of this form?

Within the definition of a class template, typename is not required when using the unqualified name of a previously declared member which declares a type; but it is required when such a reference is qualified. So the typename in "`typename A<T>::B`" is required even though "`A<T>`" is the enclosing template, "`B`" is a previously declared member, and "`B`" is recognized as a type in unqualified form. Similarly, typename is required in "`typename A::B`" since "`A`" and "`A<T>`" are synonyms within a class template with parameter list <T>.

Members of dependent base classes are never found during phase one lookup, so when they are types the typename keyword is required.

(1.17) [new issue in March]
Does an explicit instantiation directive affect the compilation model for the specified instance; for example, does it imply the “inclusion” model instead of the “separation” model, even when the export keyword is used.

The only semantics of an explicit instantiation directive are to force an instantiation (for class templates) and to add an additional point of instantiation (for function templates and members). In particular, an exported function template need not have a definition in scope at the point of an explicit instantiation directive.

(1.18) [new issue in March]
Does the “template” keyword (as applied to a dependent qualified name) apply to function and static data member templates, or just to class templates?

The "template" keyword may be applied to qualified member template names, including both class and function member templates. It may not be applied to non-template members of class templates (e.g. static data member templates), even though these entities are in some ways template-like.

(1.19) [new issue in March]
An explicit specialization declaration may not be visible during instantiation under the template compilation model rules, even though its existence must be known to perform the instantiation correctly. For example:

translation unit #1
```
template<class T> struct A { };
export template<class T> void f(T) { A<T> a; }
```

translation unit #2
```
template<class T> struct A { };
template<> struct A<int> { };  // not visible during instantiation
template<class T> void f(T);
void g() { f(1); }
```

When an explicit specialization declaration is not visible in an instantiation context, yet would affect the instantiation, the program is ill-formed but no diagnostic is required.

(1.20) [new issue in March]
According to 14.8.1, explicit template arguments may be appended to a function template name used in a call. Surely such template arguments should be allowed in other contexts in which a function name may be used, such as when taking the address of a function.

Explicit template arguments should be allowed wherever a function template name is used, not just in calls.

(2.1) [N1053 issue 6.54]
"Array/function decay in template parameter/argument lists."
The implicit "decay" of array and function types to pointer types in parameter lists, and the implicit conversion of array and function values to pointers in argument lists, should also apply to nontype template parameters and nontype template arguments.
[Core issue 758] 14 .3 [temp.arg] ¶3
"Can an array name be a template argument?"
The allowed forms for a template-argument corresponding to a non-type non-reference template-parameter do not account for the above implicit conversions; i.e. the "&" prior to an array name or function name in these cases should be optional if the values decay to pointers in the absence of "&".

As described above, from N1065. Note that this does not apply to pointers to members.

(2.2) [Core issue 759] 14.3 [temp.arg] ¶6
"Initializing a template reference parameter with an argument of a derived class type needs to be described"
It should be possible to bind a derived class object to a non-type template parameter of type reference to base class. However, the working paper only allows for standard conversions and so does not allow the derived-to-base conversions usually allowed in reference binding. This restriction should be relaxed so that such a binding is allowed.

As described above, from N1065.

(2.3) [Core issue 737] 14.6.4 [temp.dep.res]
"How can dependent names be used in member definitions that appear outside of the class template definition?"
When a member function of a class template is defined outside the class, and the return type is specified by a member of a dependent class, the typename keyword is needed to specify that the member name is a

type. So the typename keyword should be allowed in this context.

As described above, from N1065. This also applies to specifiers on static data member definitions.

(2.4) [N1053 issue 6.46]
"What are the rules used to determine whether expressions involving nontype template parameters are equivalent?"
There must be rules for determining when two template declarations/definitions refer to the same template. For template type parameters this is obvious, but when nontype parameters are used the equivalence may involve unevaluated expressions. There must be some way to determine if two such expressions are equivalent. The approach recommended in N1053 should be adopted.

As described above, from N1065.

(2.5) [N1053 issue 6.50]
"Clarification of the interaction of friend declarations and partial specializations."
[N1053 issue 8.10]
"What kind of entity can appear in a template friend declaration?"
It should be made clear that friend declarations are not allowed to declare partial specializations, and that a template friend declaration specifies that all instances of that template, regardless of whether implicitly generated and regardless of whether partially or completely (explicitly) specialized, are friends of the class containing the template friend declaration.

As described above, from N1065.

(2.6) [N1053 issue 6.53 items 1 & 2]
"Clarification of rules for partial specializations of member class templates."
When a member template of a class template is partially specialized, the partial specializations should apply to all instances generated from the enclosing class template.
When the primary template is specialized for a given instance of the enclosing class, none of the partial specializations of the original primary template should be carried over.

As described above, from N1065.

(2.7) [N1053 issue 6.58]
"Clarification of the interaction of partial specializations and using-declarations"
Using declarations only affect the visibility of declarations occurring before the using declaration itself; they do not affect the visibility of subsequent declarations with the same name. However, partial specializations of class templates are found by looking up the primary class template and then considering all partial specializations of that template. So if a using declaration names a class template, subsequent partial specializations are effectively visible because the primary template is visible. The working paper should make this clear, and should include an example.

As described above, from N1065.

(2.8) [N1053 issue 8.11]
"Clarification of conversion template instance names and using-declarations"
It should be made clear that a using-declaration (in a derived class) may not refer to an instance of a conversion function member template (in a base class).

As described above, from N1065.

(2.9) [Editorial box #8] 14.6.1 [temp.local]

The equivalence within the scope of a class template between the name of a template and the corresponding template-id should not apply when the name of the template is qualified.

As described above, from N1065.

(2.10) [Editorial box #14] 14.7.2 [temp.explicit]
An explicit instantiation directive should be a point of instantiation for each function and static data member to which the directive applies. At other points of instantiation (except end-of-translation-unit) for functions and static data members, the point of instantiation does not apply to the definition of the template unless the definition is needed at that point (e.g. inline functions, and static data members for which the the value might be required at compile time).

See 1.17

(3.1) [Core issue 780] 14 [temp] ¶1
"The definition of 'template-declaration' is incomplete"
The list of possible forms of a template-declaration does not include corresponding definitions of class members where the class is nested within a class template, nor does it include definitions of member templates (whether in non-template classes, template classes or classes nested within one of these).

As described above, from N1065.

(3.2) [Core issue 757] 14 [temp] ¶5
"Can a template member function be overloaded?"
The restriction that a function template name must be unique within a namespace scope (except for overloading) should also apply to member function templates, i.e. it should apply to class scope as well.

As described above, from N1065.

(3.3) [Core issue 781] 14.1 [temp.param] ¶8
"Must default template-arguments be written only on the first template declaration?"
The working paper should be clarified to state that default template-arguments may be specified only on the first declaration of a template in a translation unit.

As described above, from N1065.

(3.4) [Core issue 760] 14.3 [temp.arg]
"Is a template argument that is a private nested type accessible in the template instantiation context?"
It may be desirable to make it more clear (perhaps with an example) that access checking is done by name, so that if a name is accessible then it may be used in a template-id, and in the resulting instantiation there is no restriction on access to the corresponding template-parameter name itself. The working paper is already clear on this point, but has sometimes been misunderstood.

As described above, from N1065.

(3.5) [Core issue 782] [N1053 issue 6.57] 14.3 [temp.arg] ¶3
"Can a value of enumeration type be used as a template non-type argument?"
The working paper should make it clear that a constant-expression used as a template-argument for a non-type non-reference template-parameter may also have enumeration type.

As described above, from N1065.

(3.6) [Core issue 761] 14.5.1.1 [temp.mem.func] ¶3 and 14.5.2 [temp.mem] ¶3
"Can the member function of a class template be virtual?"

The term “member function template” is not used clearly here. It is not intended to mean “member template of function type”, but rather “member function of a class template which, because the enclosing class is a template, behaves somewhat like a template itself”. This distinction should be made more clear. There may be similar wording problems with respect to member templates elsewhere in the working paper.

As described above, from N1065.

(3.7) [Core issue 762] 14.5.5.1 [temp.arg] ¶4
“How can function templates be overloaded?”
An example and/or text should be added to make it clear that two distinct function templates may have identical function parameter lists and that they overload, even if overload resolution alone cannot distinguish them.

As described above, from N1065.

(3.8) [Core issue 764] 14.6 [temp.names] ¶1
“undeclared name in template definition should be an error”
In the example, the line “`T::A* a7;`” is ill-formed because “`a7`” is not dependent and has not been declared. The example should make this clear.

As described above, from N1065.

(3.9) [Core issue 766] 14.6.1 [temp.local] ¶6
“How do template parameter names interfere with names in nested namespace definitions?”
The working paper should make it clear that although class template members may hide template-parameter names, there is no such hiding with namespace members since the namespace scope is entirely outside the template declaration.

As described above, from N1065.

(3.10) [Core issue 784] 14.6.2 [temp.dep] ¶2, ¶3
“The examples in 14.6.2 on dependent names need work”
[Public comments #7, #23]
Some of the examples in this section are in disagreement with the textual description of dependent names and lookup rules. The examples should be corrected or removed.

As described above, from N1065.

(3.11) [Core issue 767] 14.6.4.1 [temp.point]
“Where should the point of instantiation of class templates be discussed?”
There should be cross-references between the various paragraphs discussing points of instantiation, with respect to class, function and static data member templates.

As described above, from N1065.

(3.12) [Core issue 786] 14.7.2 [temp.explicit]
“The description of explicit instantiation does not allow the explicit instantiation of members of class templates (including member functions and static data members)”
The description should be extended to include all the members, and members of members, for which explicit instantiation is appropriate.

As described above, from N1065.

(3.13) [Core issue 787] 14.7.3 [temp.expl.spec]
"Make it clear that a user must provide a definition for an explicitly specialized template; if not, the program is ill-formed"
It should be clear that when a template is explicitly specialized, the unspecialized template is not used and so there is no implicit generation for the specialization. Therefore if the specialization is used it must be defined, following the rules for requiring definitions for non-template declarations. (In particular, there are some cases where a diagnostic is required and some where no diagnostic is required.)

As described above, from N1065.

(3.14) [Core issue 677] 14.8.2 [temp.deduct]
"Should the text on argument deduction be moved to a subclause discussing both function templates and class template partial specializations?"
There should be cross-references between the various places where template argument deduction is done.

As described above, from N1065.

(3.15) [Core issue 768] 14.8.2 [temp.deduct] ¶10
"typename keyword missing in some examples"
In the specified paragraph, the typename keyword is requires in two places:

```
T deduce(typename A<T>::X x,   // T is not deduced here
         ^^^^^^^^
         T t,                  // but T is deduced here
         typename B<i>::Y y);  // i is not deduced here
         ^^^^^^^^
```

As described above, from N1065.

(3.16) [Core issue 788] 15.3 [except.handle] ¶9
"Is it implementation defined whether the stack is unwound before calling terminate in all of the 8 situations described in 15.5.1?"
It should be made clear that in all other cases where terminate is called (other than due to failure to find a matching handler), the stack is not unwound. Also, there are other cases where an implementation might determine, before finishing a stack unwind, that terminate will be called during the unwind. The working paper should specify whether that portion of the unwind must actually be done.

As described above. Also,while handling an exception, an implementation may not terminate stack unwinding early based on a determination that the unwind process will eventually fail and cause a call to terminate().

(3.17) [Editorial box #13] 14.7 [temp.spec] ¶1
This paragraph does not really describe the handling of member templates and of members of classes nested within class templates. The missing cases should be added.

As described above, from N1065.

(3.18) [Public comment #6]
It should be made clear that a class template is instantiated in any context where the completeness of the type might have an effect on the semantics of the program.

As described above, from N1065. Note that overload resolution is a special case which has already been dealt with.

(3.19) [Public comment #8] 14.6.2 ¶5

The phrase “If a template-argument  is a used as a base class...” should be changed to match the intent in ¶4, e.g. “If a base class is a dependent type...”.

As described above, from N1065.

(3.20) [Public comment #12]  14.5.3 ¶4
The phrase “the corresponding member function” is incorrect; the friend might be a class.  So the word “function” should be deleted.

As described above, from N1065.

(3.21) [Public comment #20]  15
Clause 15 should explicitly state that multiple exceptions may be active at the same time (“recursive” exceptions).  The current wording implies this but never explicitly says that this is allowed.

As described above, from N1065.

(3.22) [Public comment #20] 15.1 ¶1
The example needs to be updated to account for the new type of string literals.  Also it might be useful to point out that the special implicit cv-qualification conversion for string literals does not apply to throw-expressions.

As described above, from N1065.

(3.23) [Public comment #23]  14.7.2, 14.7.3
The situations in which an empty template argument list “<>” may be omitted should be more clearly explained, particularly in the examples in these sections.

As described above, from N1065.

(3.24) [Public comment #23] 14.7.3 ¶16
The phrase “A member template ... is not be implicitly ...” should read “A member ... is not implicitly ...”.

As described above, from N1065.

(3.25) [Public comment #23] 18.6.2.2 ¶2
The description of “unexpected” differs from the description in 15.5.2.The description in 15.5.2 is correct; the one in 18.6.2.2 should either be changed to match or be replaced with a cross-reference to 15.5.2.

As described above, from N1065.

(3.26) [Public comment #24]  15.1 ¶2
The wording in this paragraph about exiting a try block should actually refer to exiting just the “try” portion of the try construct.  That is, a throw from within a handler should never be caught by that handler or by a handler associated with the same try.

As described above, from N1065.

(3.27) [Public comment #28]  14.7.3 ¶6, ¶16
The examples in these two paragraphs contradict each other.  It appears that the last line of the example in paragraph 16 should not contain “<>” because the definition should not be an explicit specialization.

As described above, from N1065.

(3.28) [Public comment #29]
The semantics, use and intent of the keyword “export” need to be clarified.

As described above, from N1065.

(3.29) [Public comment - not yet numbered]
The working paper has rules for handling a “>” within an expression in a template-id (14.2 ¶3). A similar ambiguity occurs with expressions written as default arguments for nontype template parameters in the parameter list of a template. The same solution should apply.

As described above, from N1065.

(3.30) [new issue in March] 14.6.2 [temp.dep] ¶4
The sentence “`X<T>::a` has type `double`.” should be moved to a comment in the example, as in:

```
template<class T> struct X : B<T> {
        A a;   // "a" has type "double"
};
```

As described above, from N1065.

(3.31) [new issue in March]
The working paper should explicitly state that the implicit instantiation of a class template does not cause the implicit instantiation of the definition of a static data member, and therefore does not (by itself) cause the initialization (and associated side-effects) of static data members to occur.

As described above, from N1065.

(3.32) [new issue in March]
The working paper should explicitly state that an entity which appears to be “used” in a default argument is actually used only if the default argument itself is used.

As described above, from N1065. Note that this effects both template instantiation and the requirement that a definition exist if it is used, even for nontemplates.

(3.33) [new issue in March] 2.3.1.2 ¶3
They keyword “friend” on the first line of the example should be removed.

As described above, from N1065.

(3.34) [new issue in March] 14.7.3 [temp.expl.spec] ¶3
It should be made clear that the restriction on default arguments “in” explicit specializations applies only to function template explicit specializations (including member functions and member function templates where the enclosing class is not specialized), and not to member functions of class template specializations (which are not themselves specializations).

As described above, from N1065.

(3.35) [new issue in March]
The restrictions (in 8.3.6 ¶4 and 8.3.6 ¶6) on default arguments in templates are not sufficiently complete; for example, they do not specifically mention member functions of class templates and member templates.

As described above, from N1065.

New issues not discussed in N1065

New issue (core issue 771)

Two functions with the same name and different return types, but otherwise identical, may not be declared in the same namespace, even if they are in different translation units. Does the same restriction apply to function templates?

No change. The draft currently allows this.

New issue (see core reflector 7399)

When is the "export" keyword permitted, and when is it required?

- The export keyword may appear only on template declarations at namespace scope and template definitions at namespace scope.
- A template is exported if declared with "export" on any declaration or definition.
- A template must be defined in any translation unit in which it is declared with the "export" keyword.
- A template may not be declared export for the first time after its definition.

This is minimum change to the working paper which appears to have significant support.

New issue

What standard conversions can apply to a template argument to bring it to the type of the corresponding nontype template parameter?

For parameters of integral or enumeration type, only the integral promotions and integral conversions are allowed (and not, for example, floating/integral conversions). For pointer and reference parameters, only derived-to-base conversions and conversion to "`void*`" are allowed. (If "`void&`" is added and conversion to "`void&`" is a standard conversion, then this would be allowed also.)

(Note that array-to-pointer and function-to-pointer conversions would always be done under the proposed resolution to 2.1; see 2.2 also.)

(Note that for pointers and references, the access checking is done at the template-id but the conversion is done within the template; for example, consider conversions across virtual inheritance.)

New issue

Does a friend declaration ever refer to a namespace-scope template declaration with the same name?

No, except that when explicit template arguments are provided the friend declaration refers to the specialization. Other than that, if the enclosing scope has a function template and no non-template function, the friend declares a (hidden) new nontemplate function, exactly as if the function template declaration did not exist.

New issue

The "`typename`" keyword is not permitted in a function-style cast, e.g. "`return typename T::X(y);`".

Change the grammar to allow this case.

Core issue 783

A name in a dependent base class may hide a template parameter name when a member function of a class template is instantiated.

No change. This is similar to many other troublesome hiding problems in C++. It is too late to address such issues now, although an implementation might warn about this kind of hiding.

New issue

The example in 14.6.3 is unclear.

The example in 14.6.3 should make it clear that although an implementation is allowed to diagnose this kind of error when processing the template definition, it is not required to diagnose such errors until the point of instantiation.

New issue

In the following example:

```
namespace A {
    struct B { };
    template<class T> void f(T t);
}
void g(A::B b) {
    f<3>(b);
}
```

does type-dependent (Koenig) lookup apply to the lookup of "f"? Without the explicit "<>" template arguments, the answer is currently yes, because the lookup of "f" (other than to determine whether it is a type) can be deferred until after the arguments have been parsed. But with explicit template arguments, there is no way to parse without knowing that "f" is a template.

The "template" keyword should be allowed in this context so that type-dependent lookup can be used even when there are explicit template arguments.

**Unresolved Issues**

From N1065, the Core-III list of issues compiled at Nashua

(1.1) [Core issue 763] 14.5.5.2 [temp.func.order] ¶2 and 14.5.2 [temp.mem]
"Partial Specialization: the transformation also affects the function return type"
Partial ordering of function templates is supposed to be used to help disambiguate in the selection of the best conversion function template, but the mechanism described for establishing a partial ordering between function templates does not work for conversion function templates.

The rules for partial ordering of member template conversion functions should be based on a transformation to an ordinary function. This needs work - it is not as simple as it sounds - and John Spicer will investigate further for the July meeting.

(1.15) [new issue in March]
The partial ordering rules for function templates are overly restrictive: they require that two functions being compared have identical signatures (13.3.3 ¶1). This restriction could be relaxed to just require that the functions have identical parameter types for overloading purposes.

The rules for overloading and partial ordering of function templates can produce some surprising results due to the fact that the cv-qualification and lvalue-to-rvalue conversions are used to eliminate some functions before the partial ordering rules are applied.

The majority of the informal core-III group preferred no change here, but John Spicer strongly disagreed. John Spicer and Steve Adamczyk wrote a paper on this issue for the mailing; the document numbers are 97-0046 / N1084.

(1.16) [new issue in March]
The working paper allows member template conversion functions, and implies that their template parameters may be deduced, but does not specify the deduction rules. These rules must be stated explicitly.

The rules for template argument deduction for member template conversion functions must be written. John Spicer volunteered to investigate this issue.

New issues not discussed in N1065

New issue
What are the deduction rules for the use of template functions in contexts other than a function call, e.g. taking the address of a function template (yielding the address of a particular specialization)? See issue 1.8 above.

New issue
A "constant" nontype template parameter may be used as a template argument, even though it is not actually a constant at its point of declaration. Similarly, it should be possible to use a pointer or reference nontype template parameter as a template argument even though it is not actually the name or address of an external-linkage object at its point of declaration.

**Recommendations**

The committee members attending the informal core-III issues meeting agreed, in most cases unanimously, to recommend that the full core-III group and the full committee approve the above proposed issue resolutions at the July meeting.

Of the unresolved issues, only 1.15 is expected to require extensive discussion. The other issues are expected to be non-controversial once the technical details are worked out.

The informal group also discussed template template parameters; see document 97-0043 / N1081.